



Create a 20 Times Faster Database Engine Optimized to MMOGs

Shuichi Kurabayashi

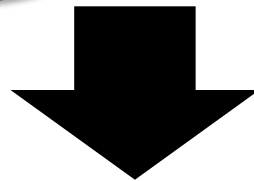
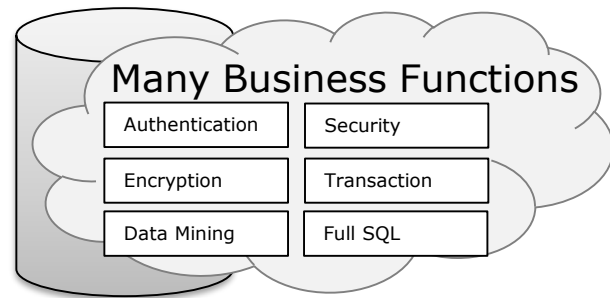
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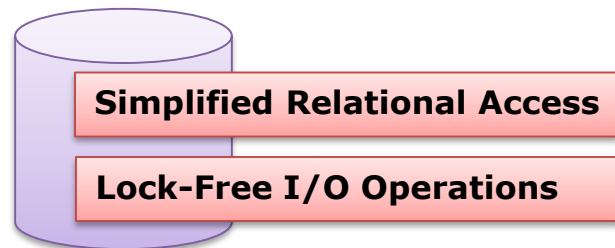
Summary

A **simple in-house database engine tailored to MMOGs** is highly effective to provide mobile MMOGs for multi-millions unique users.

Legacy Full-Featured DBMS



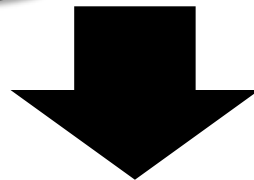
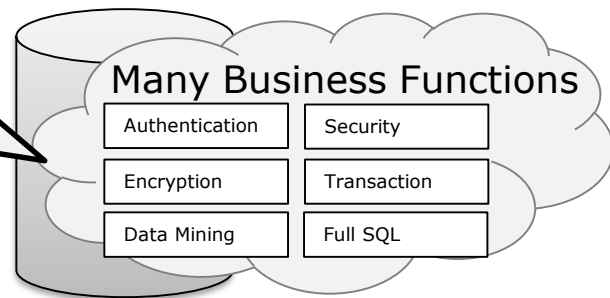
Minimal & Tailored DBMS



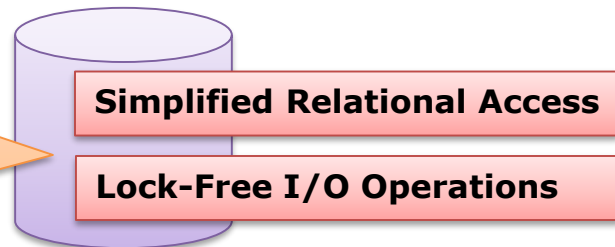
Game developers do not use almost of those business-oriented functions

It is easy to implement the **minimal** database management system.

Legacy Full-Featured DBMS



Minimal & Tailored DBMS



My message is:

**Let's create your own
database engine
tailored to your titles!**

Quick Introduction

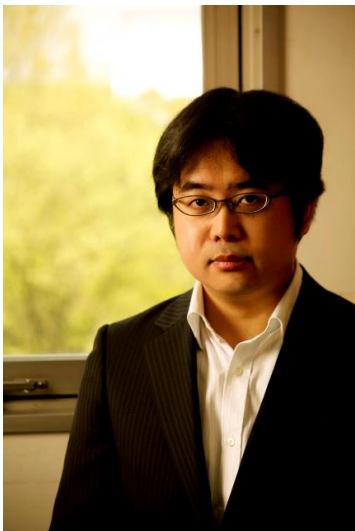
Who are you?

What is Cygames?

What is problem?

Introduction: Who I am

Shuichi Kurabayashi, Ph.D.



- Technical advisor of Cygames, Inc. Also Director of Cygames Research.
- Project Associate Professor at the Graduate School of Keio University

Introduction: Cygames is one of the largest mobile game developers in Japan

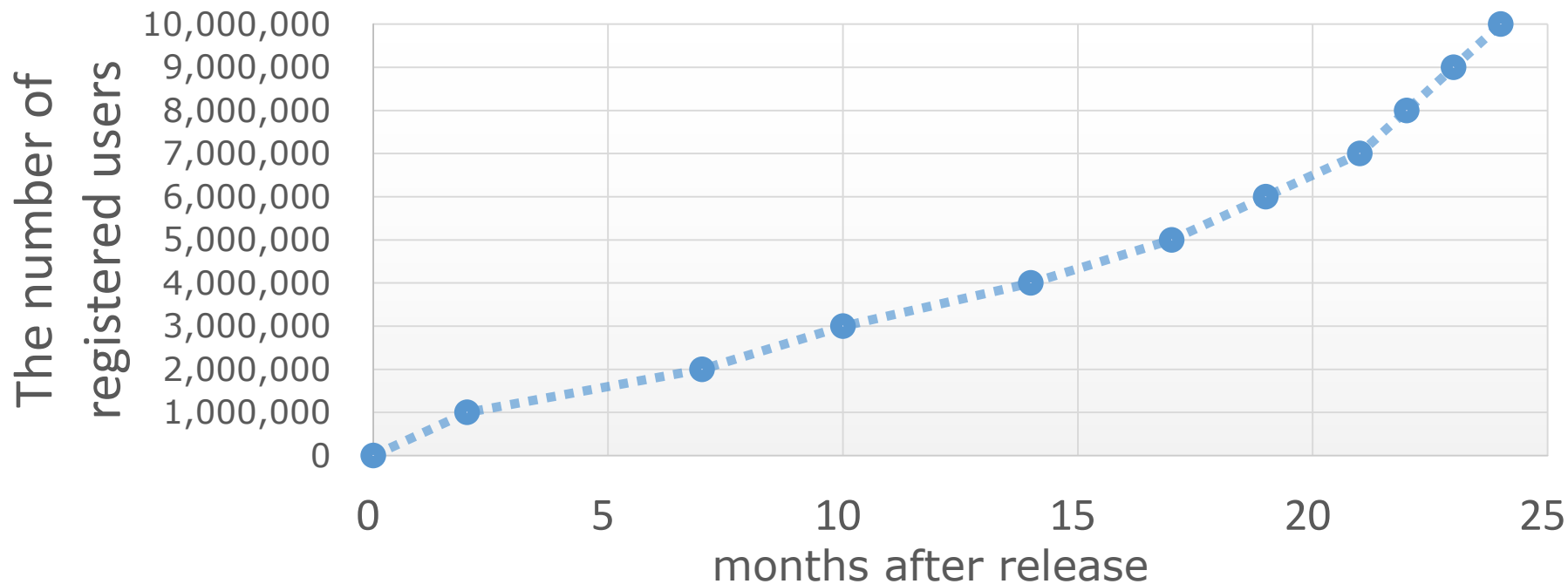


Known as the developer of the wildly popular card battle game "Rage of Bahamut".



Recently released "Shadowverse".

Background: We have been providing mobile MMOGs for 10 millions users.



Background: Mobile gaming is an important subject of DB Research

Tokyo Stock
Exchange

6,700/sec
(400,000/min)
transactions

Twitter

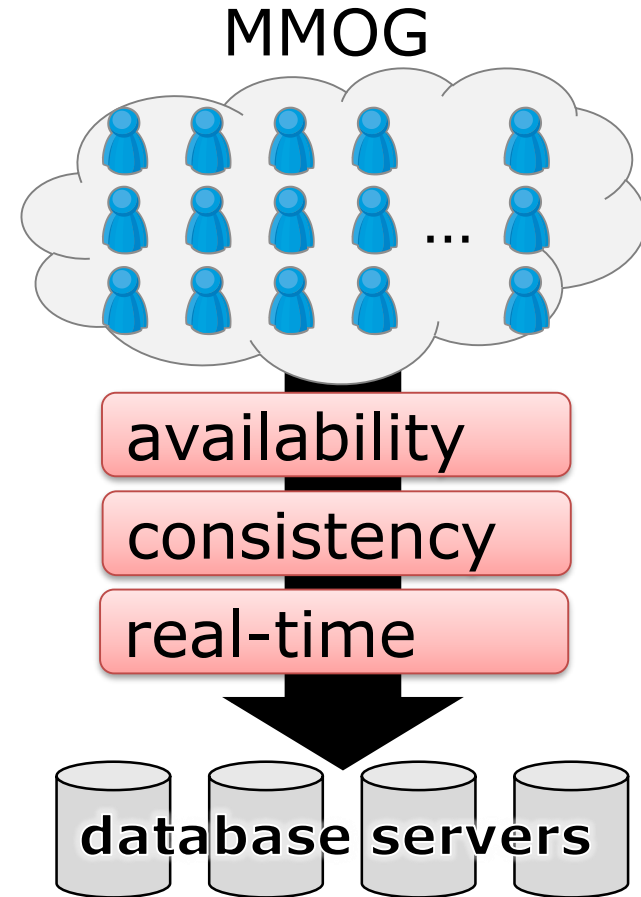
40,000/sec
messages at
peak.

**Mobile game
(Japan)**

**100,000/sec
transactions
in average.**

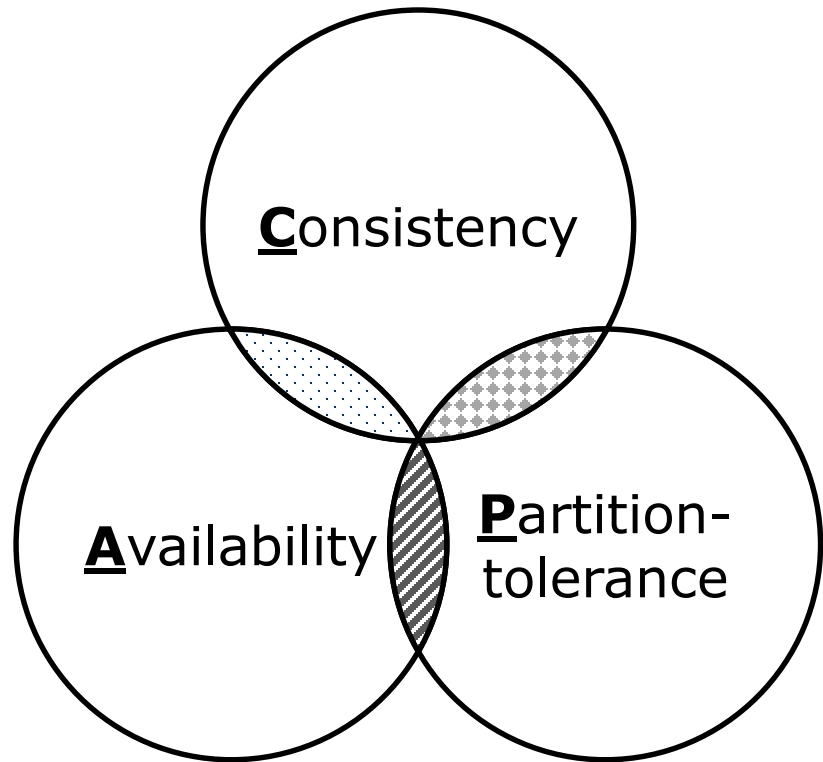
Problem Definition

- Modern MMOGs require databases to support large-scale **availability**, strong **consistency**, and **real-time** response.
- It is difficult to support such capabilities efficiently by using conventional systems.



CAP Theorem

- **C: Consistency of data**
 - After data has been updated, if something else references that data, it will always be guaranteed that the updated data can be referenced.
- **A: Availability of the system**
 - No matter what the current situation, the system will continue to operate.
- **P: Tolerance to network partitions**
 - Data can be distributed.



CAP Theorem

- **C: Consistency of data**

- After data has been updated, if something else references that data, it will always be guaranteed that the updated data can be referenced.

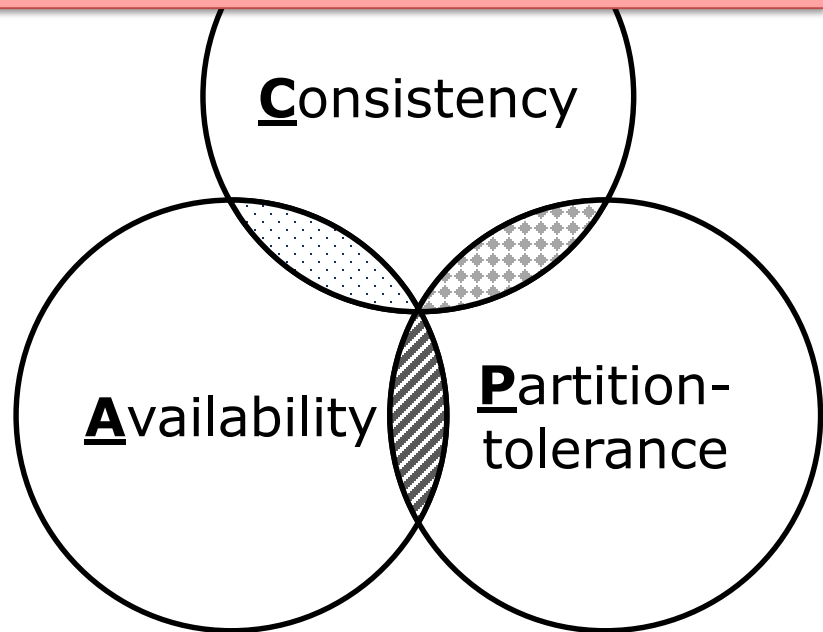
- **A: Availability of the system**

- No matter what the current situation, the system will continue to operate.

- **P: Tolerance to network partitions**

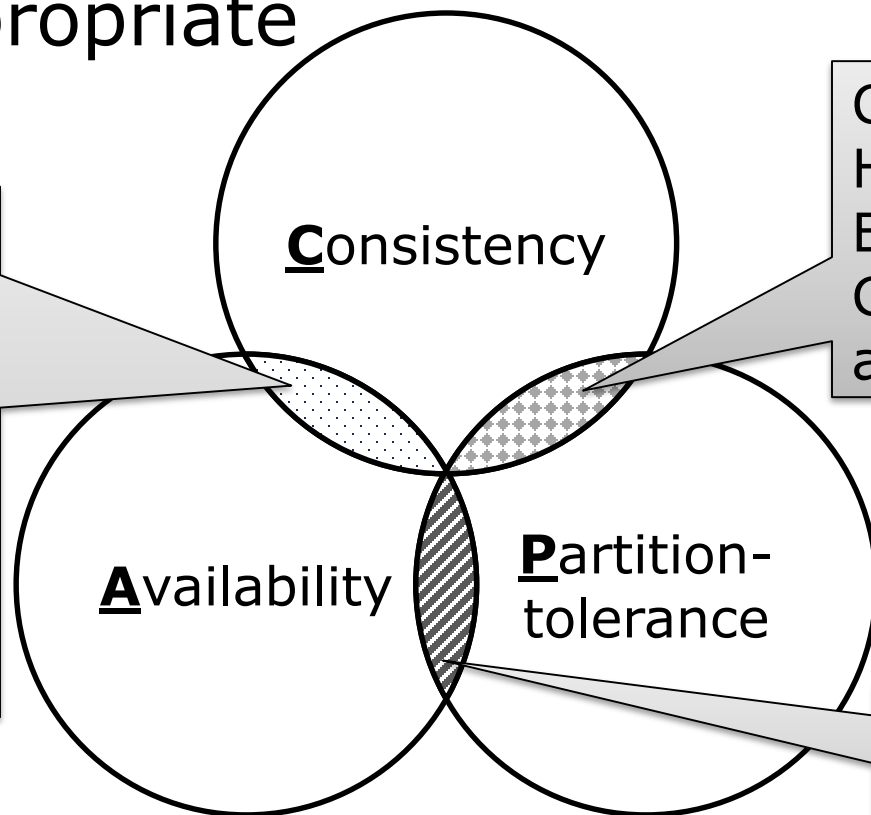
- Data can be distributed.

Out of these three guarantees, only two can be fulfilled at a time.



Choosing Appropriate System

CA: conventional relational database systems, such as MySQL, Oracle, and SQL server, strong consistency

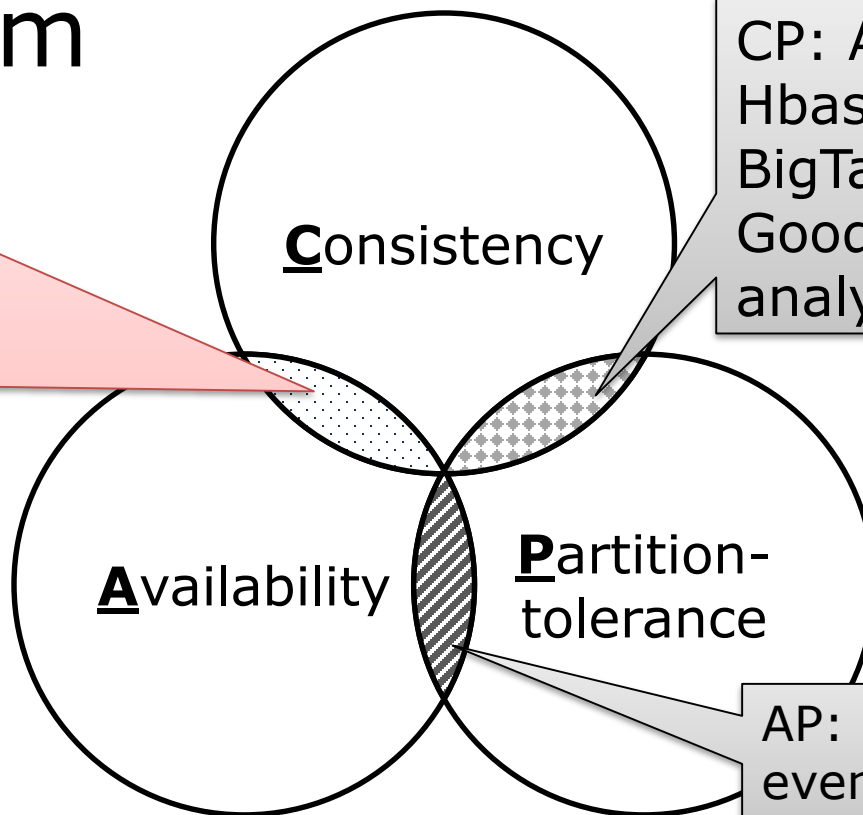


CP: Apache Hbase, Hadoop, BigTable, Good for analytics

AP: SNS, NoSQL, eventual consistency

CAP Theorem

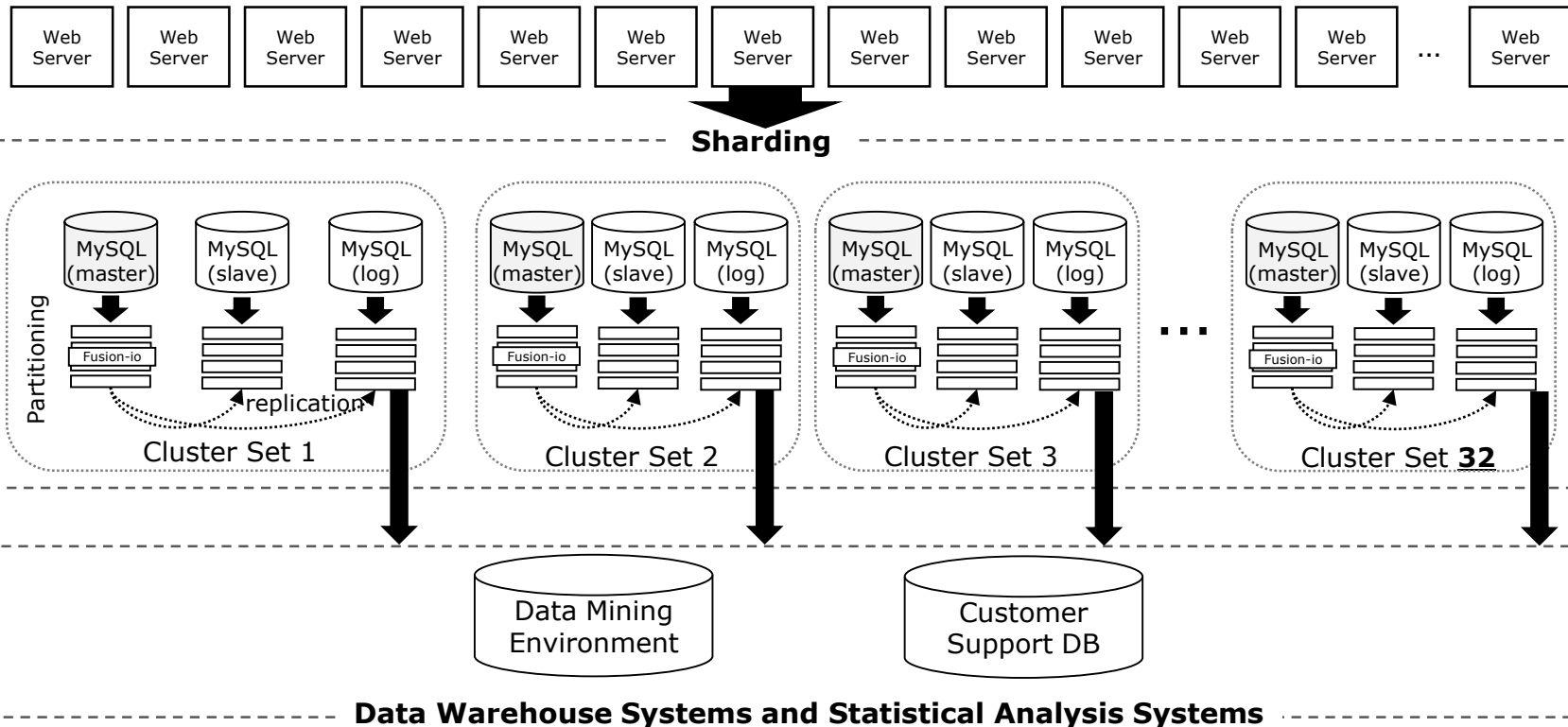
CA is the most important for games, because games require **strong consistency and fast response**



CP: Apache Hbase, Hadoop, BigTable, Good for analytics

AP: SNS, NoSQL, eventually consistency

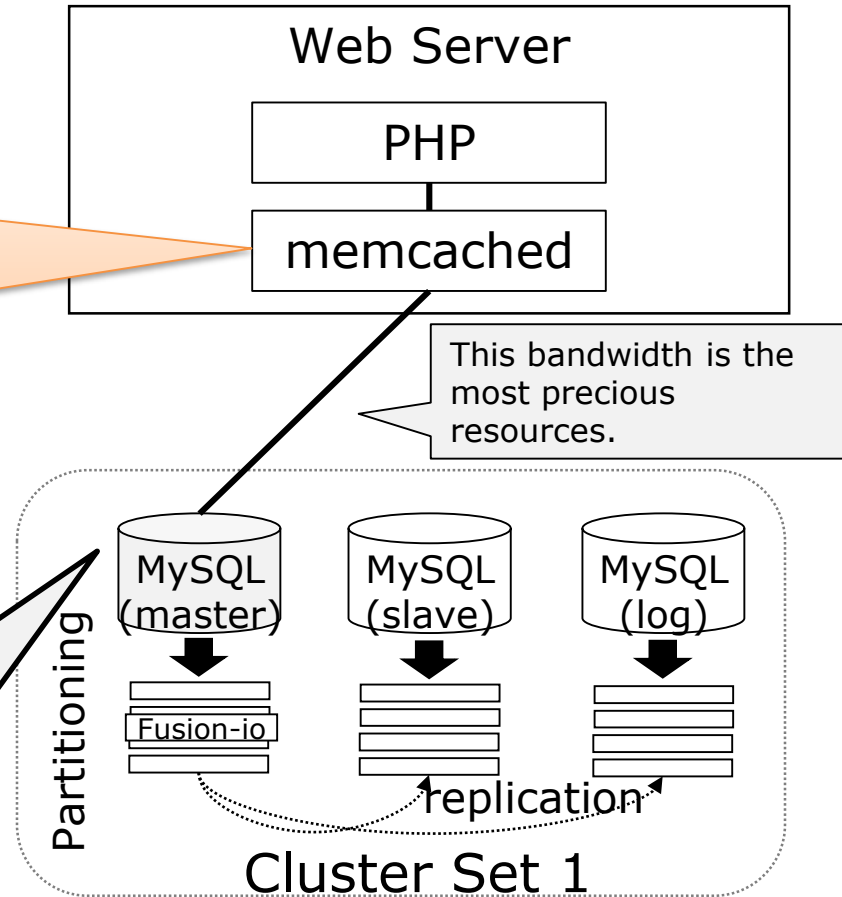
Typical Backend Architecture



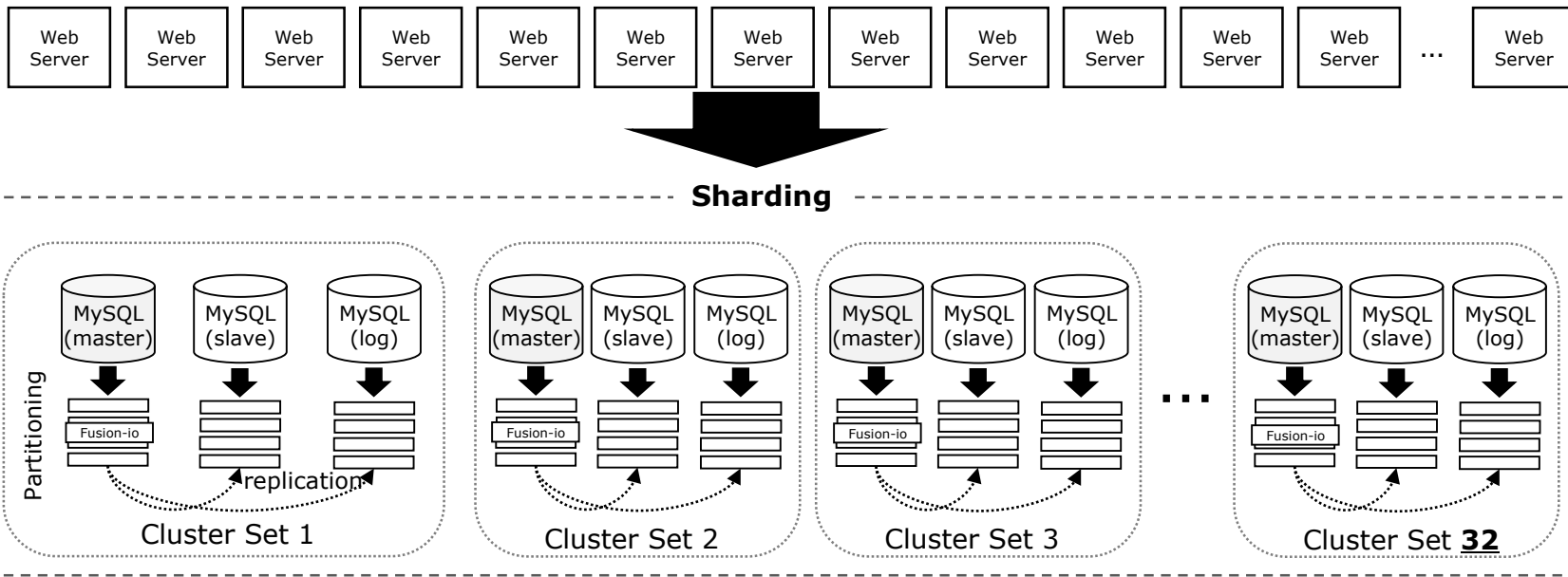
Primitive Unit

Main memory, memcached stores static data shared among all the PHP processes.

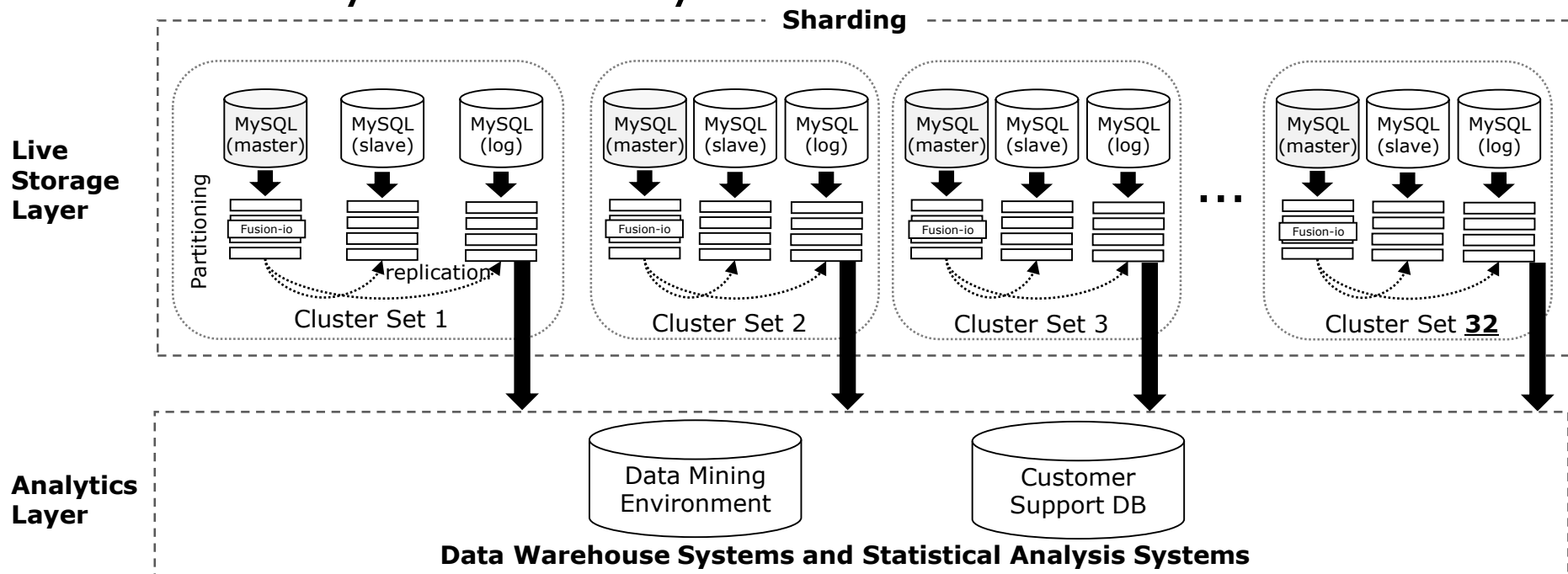
MySQL stores status data, which are updated in a real-time manner.



Combination of **sharding** (horizontal decomposition) and **partitioning** (vertical decomposition) brings high-level parallelism



Data mining process including JOIN operations is carried out by replica DBs that are replicated from the master database asynchronously.



MySQL is working now. But...

Operation can be highly inefficient and unprofitable

Not so
scalable

Large
maintenance
cost

Capacity per
machine is
not so high

RDB functionalities are appropriate for C-A requirements, but it is too fat.

3 Overheads that decrease RDB's performance

Overhead-1: Too Generic Structure

Overhead-2: Not utilizing data access pattern in games

Overhead-3: Not utilizing modern hardware

Overhead-1: Too Generic Structure!

Data Model
supported
by Apps

External
Schema-1

External
Schema-2

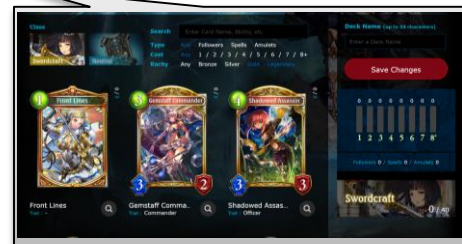
External
Schema- n

Data Model
supported by DBMS

Conceptual Schema

Data Model
supported by storage

Internal Schema

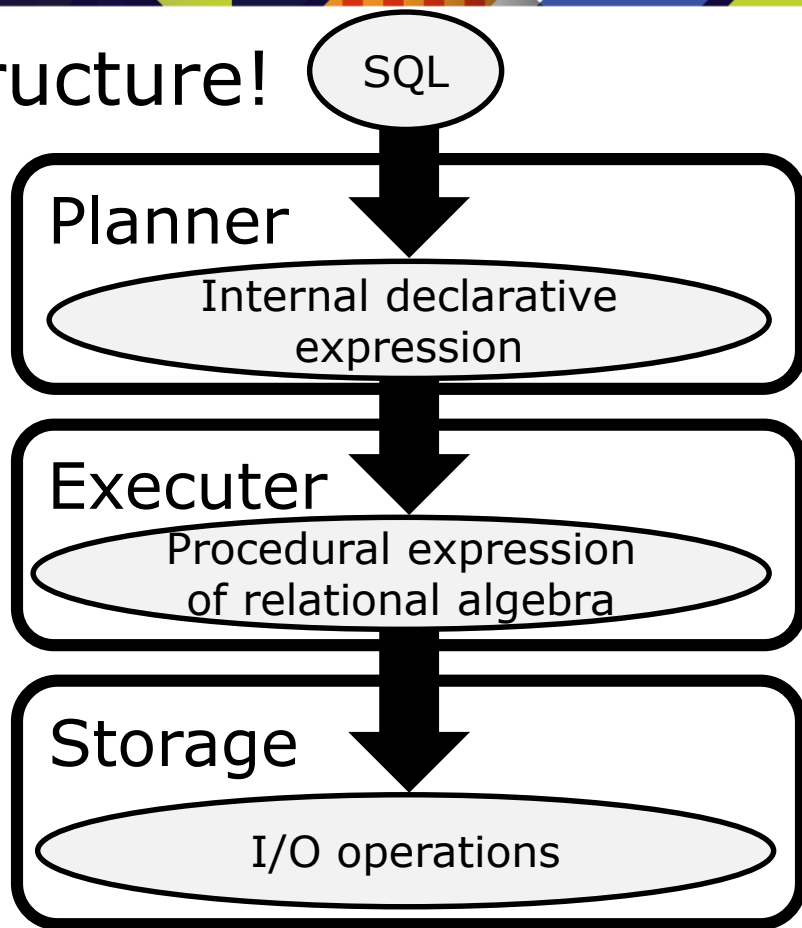


`CREATE TABLE (...);`

`TupleTableSlot *rec;`

Overhead-1: Too Generic Structure!

- Planner (logical optimization)
 - Rewrites queries for better performance, by analyzing queries as relational calculus.
- Executor (physical optimization)
 - Executes query primitives such as relational algebras.
- Storage Subsystem
 - Reads and writes disk storages.



Let's resolve Overhead-2: Not utilizing data access pattern in games

	Insert	Update	Select	Real-time	Consistency	Query Type
Legacy Web (e-commerce)	Small	Small	Large	No	Strong Consistency	Dynamic (mutable)
SNS(Large Scale Web Apps)	Large	Small	Massively Large	No	Eventual Consistency	Dynamic (mutable)
IoT (Stream DB)	Massively Large	no	Small	Yes	Application-Dependent	Static (immutable)
Mobile Game	Small	Massively Large	Massively Large	Yes	Strong Consistency	Static (immutable)

We faces technical requirements that are essentially different from a conventional DBs

Let's resolve Overhead-3: Not utilizing modern hardware

The throughput performance becomes from 10 to 20 times faster, due to the multi-core CPU with the high parallelism and a SSD with the highly parallel access.

Solution

Query Pre-Compilation
(Immutable Query)

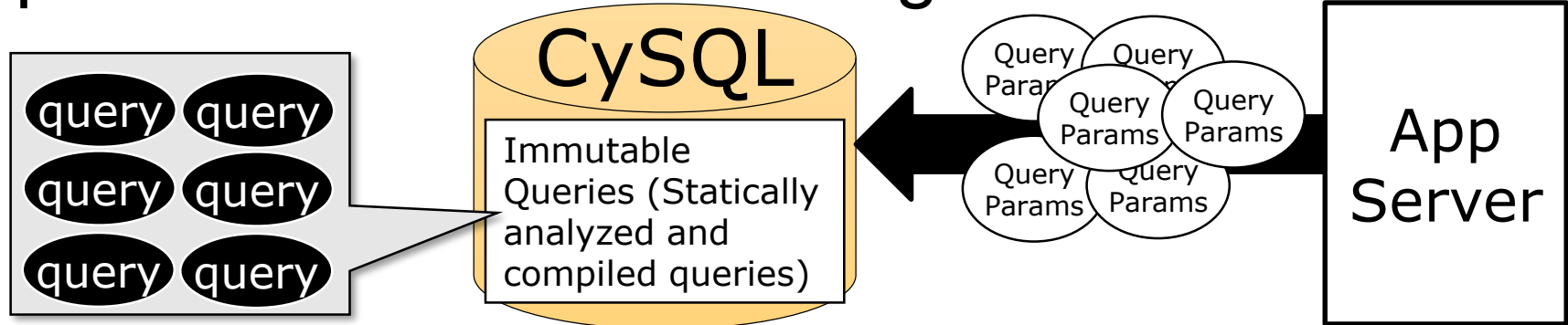
```
graph TD; A[Query Pre-Compilation (Immutable Query)] --> B[Lower Footprint query processing]; A --> C[Lock-Free Thread Scheduling];
```

Lower Footprint
query processing

Lock-Free Thread
Scheduling

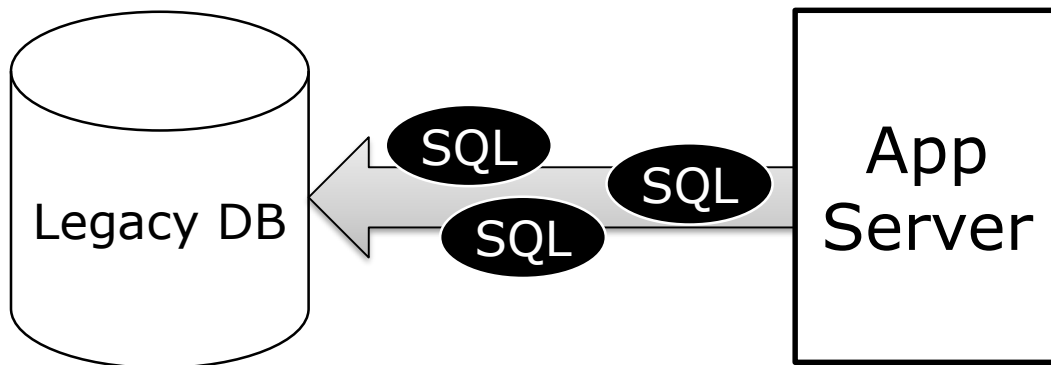
Static Query Analysis is effective

If we can know all the query to be processed, we can apply holistic optimization by analyzing their data access pattern and implicit race condition among them.



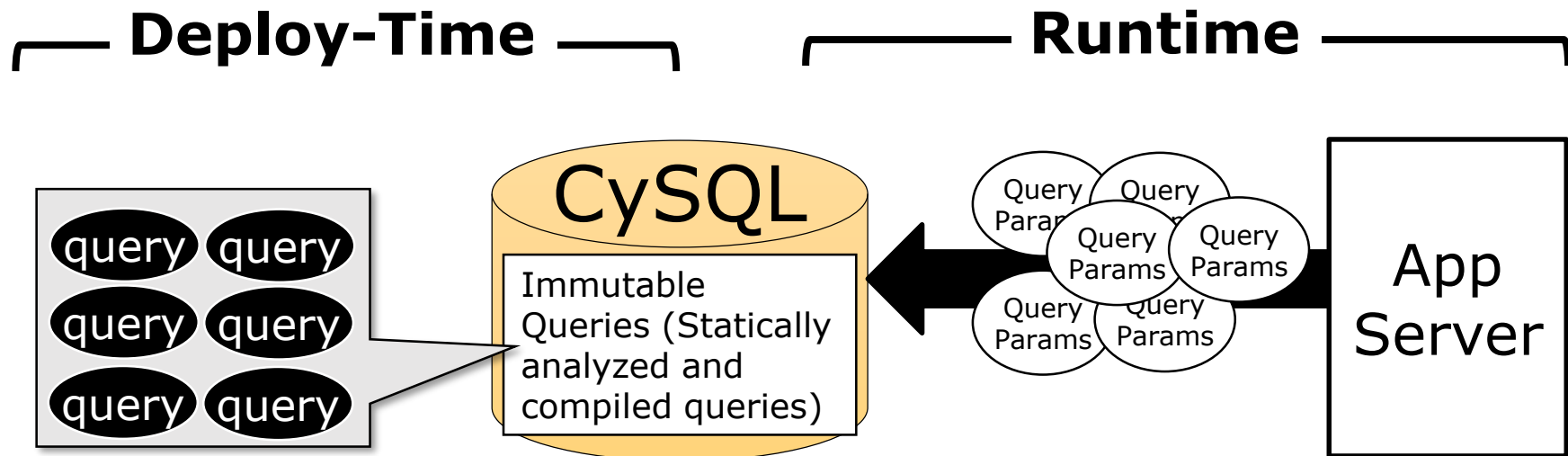
Legacy Query Model

- Legacy query model cannot optimize holistic query, because queries are submitted dynamically.



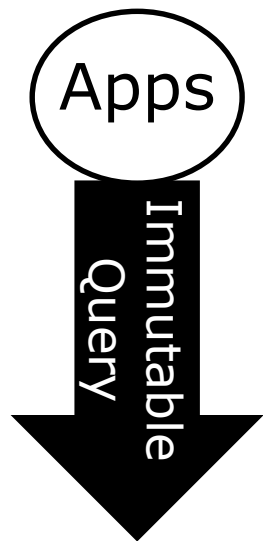
Immutable Query

Immutable query model requires the all queries to be defined at deploy time



SQL parsing and optimizing takes large cost

Controlling transaction manager directly from the application, we can achieve nearly 10 times faster performance than SQL



Planner

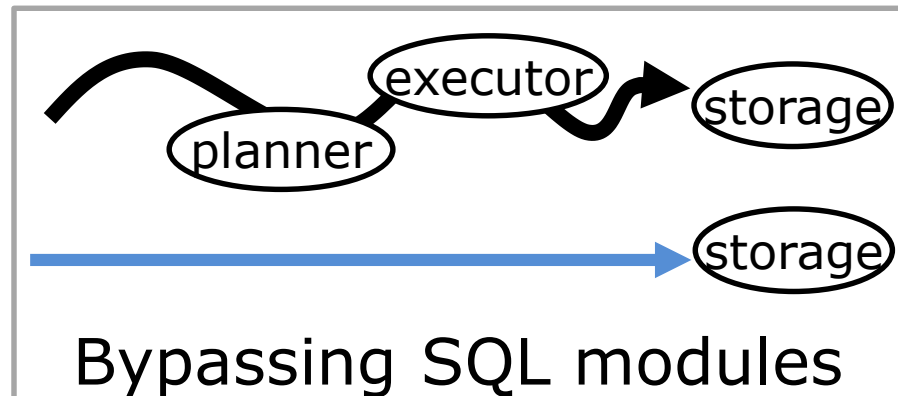
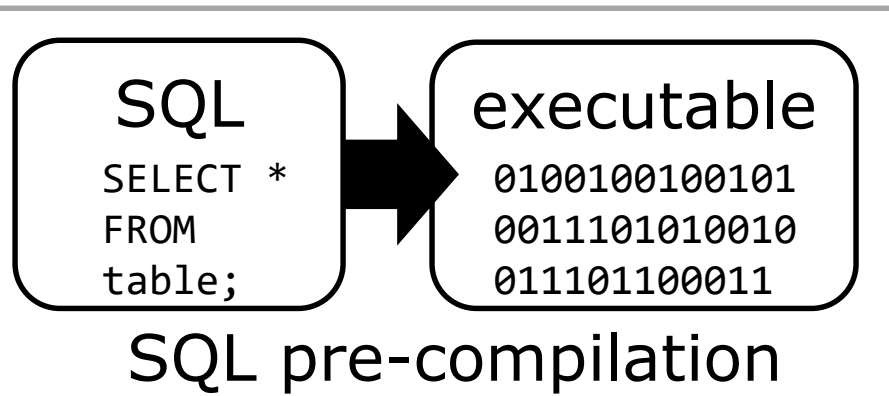
Executer

We can bypass the planner and the executor, by compiling SQL into executable machine code.

Storage

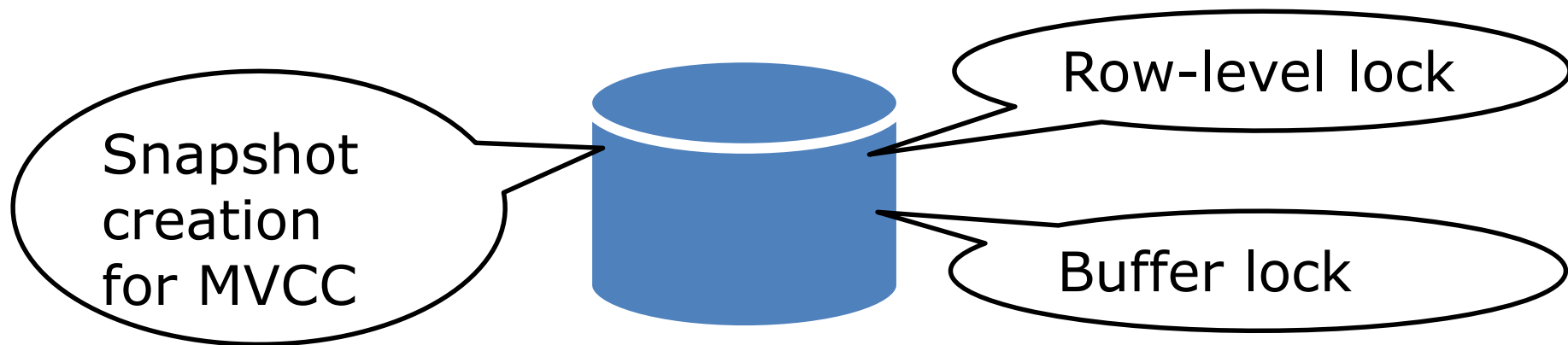
Query Processor Bypassing

Introducing query pre-compilation and bypassing planner and executor



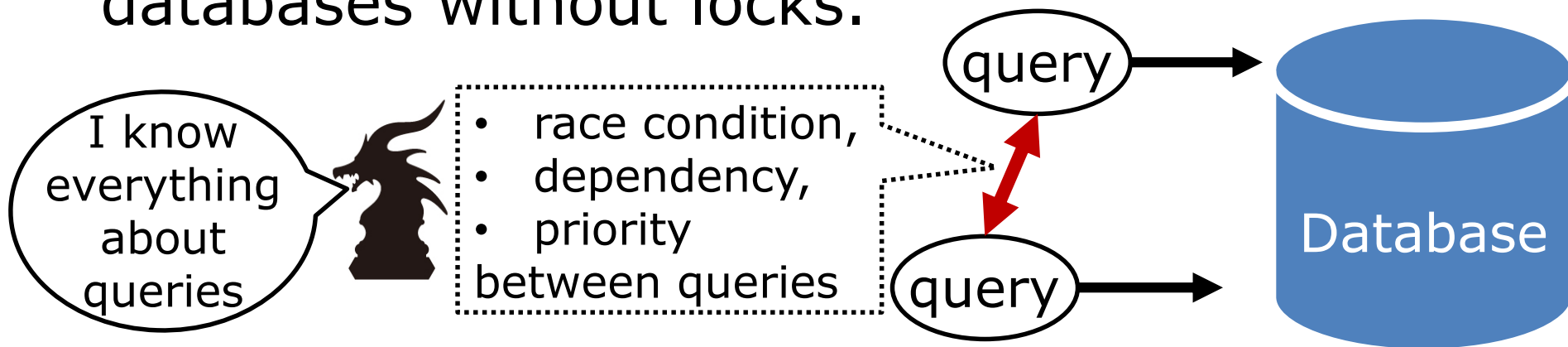
Lock-Free Transaction Processing

Legacy RDB embraces a lot of overhead of sync inside the DB engine for lock mechanism. Each of those overheads is small, but small overhead will be accumulated, so each overhead is the cost which can't be ignored.

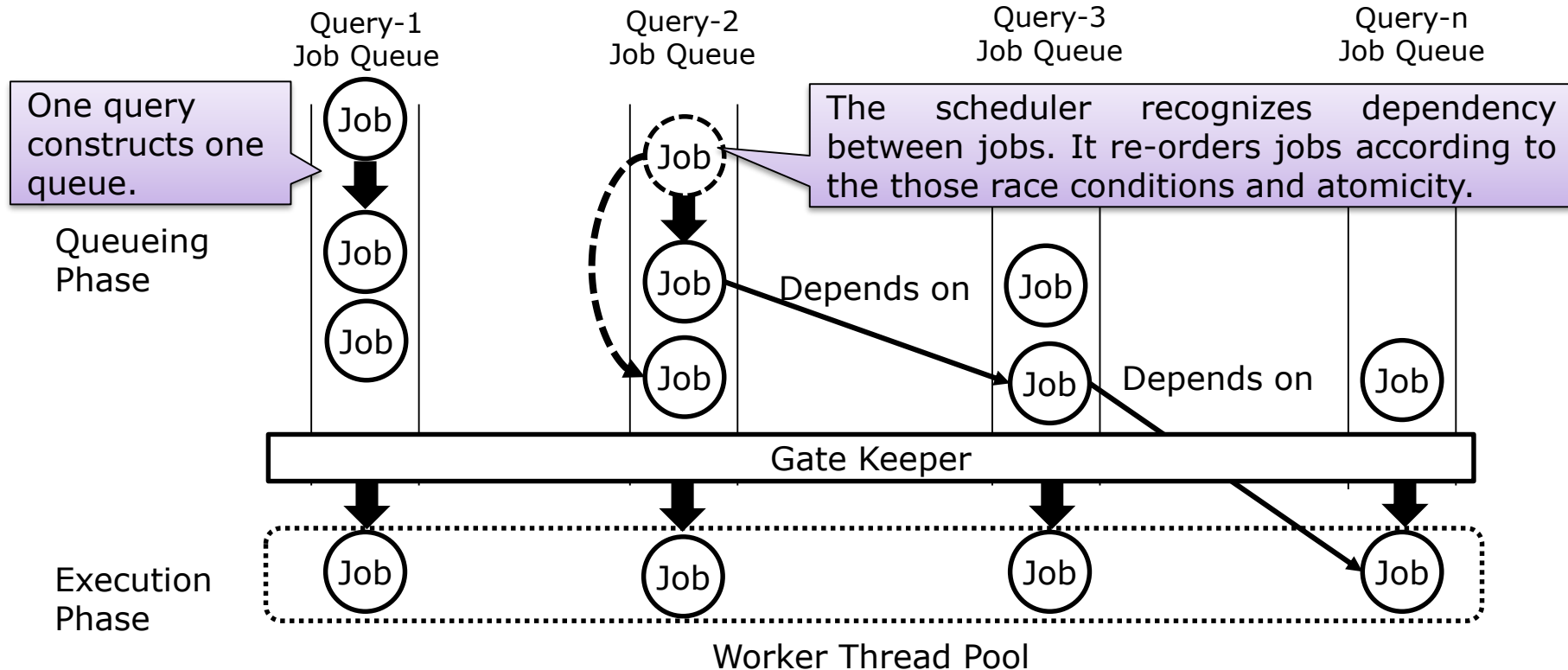


Can we realize Lock-Free?

By using the thread scheduling specialized in the data access pattern specific to the game, multiple threads read/writes databases without locks.



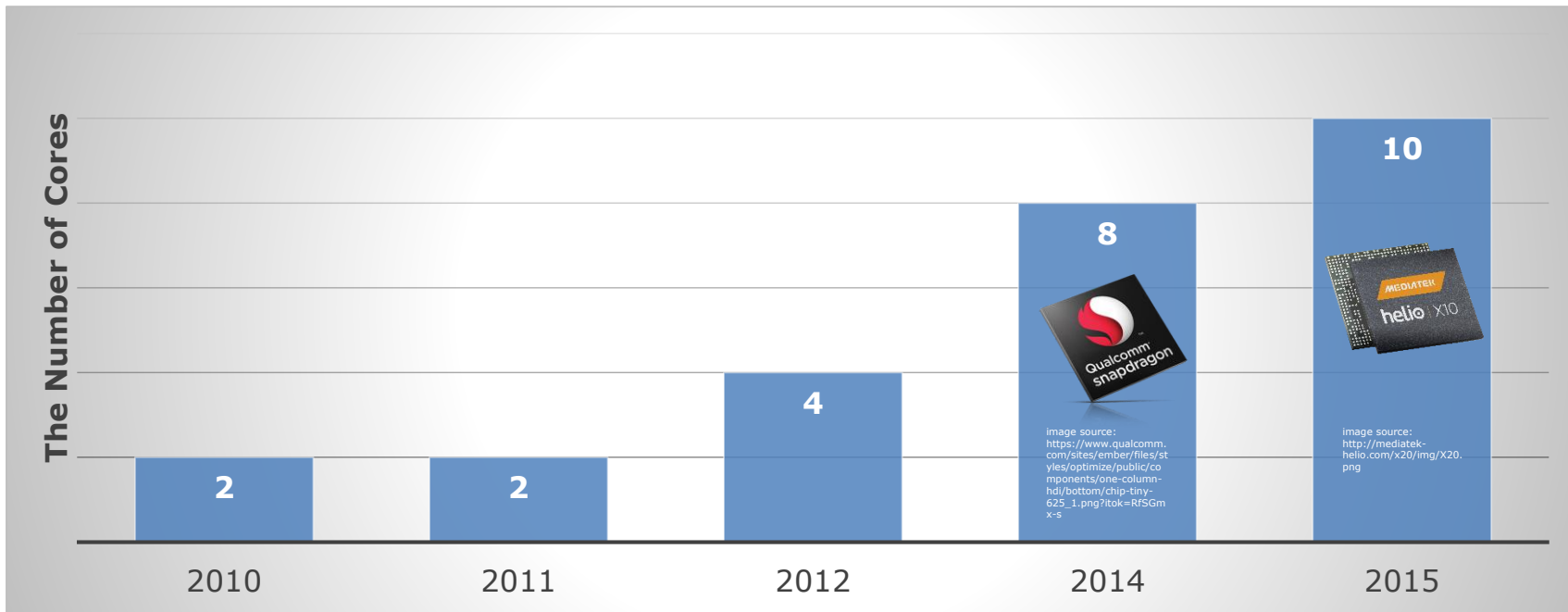
Basic Concept of Lock-Free Thread Scheduling



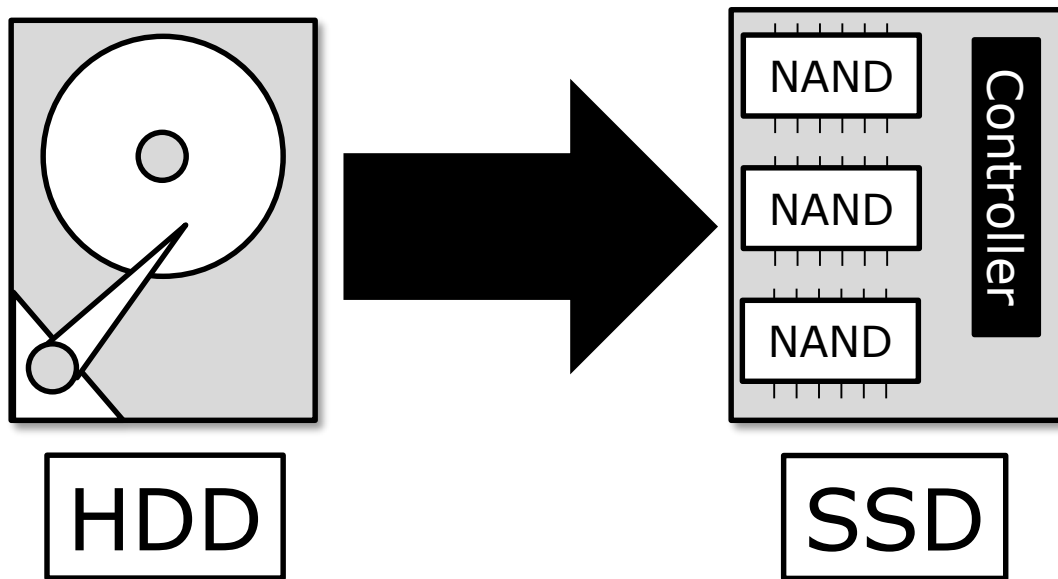
In the immutable query model, DB engine can schedule threads perfectly without synchronization lock.

- Because the DB engine can analyze race condition between queries statically.
- The thread scheduling strategy is easy: "Do not schedule the queries that cannot be executed simultaneously."

The number of CPU cores are increasing for both energy and performance efficiency.



Parallelism in I/O Devices

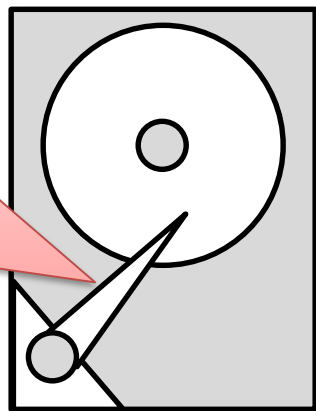


Not so parallel

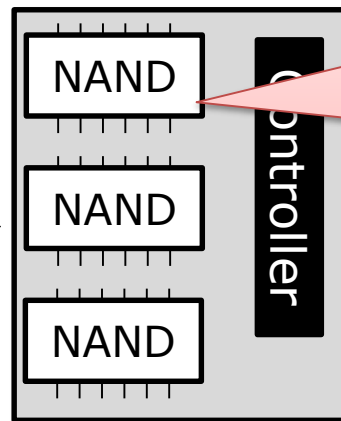
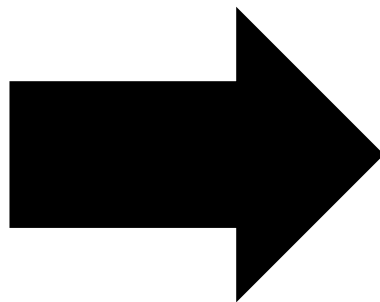
Highly Parallel

Parallelism in I/O Devices

Only this parts (head) can read/write data



HDD



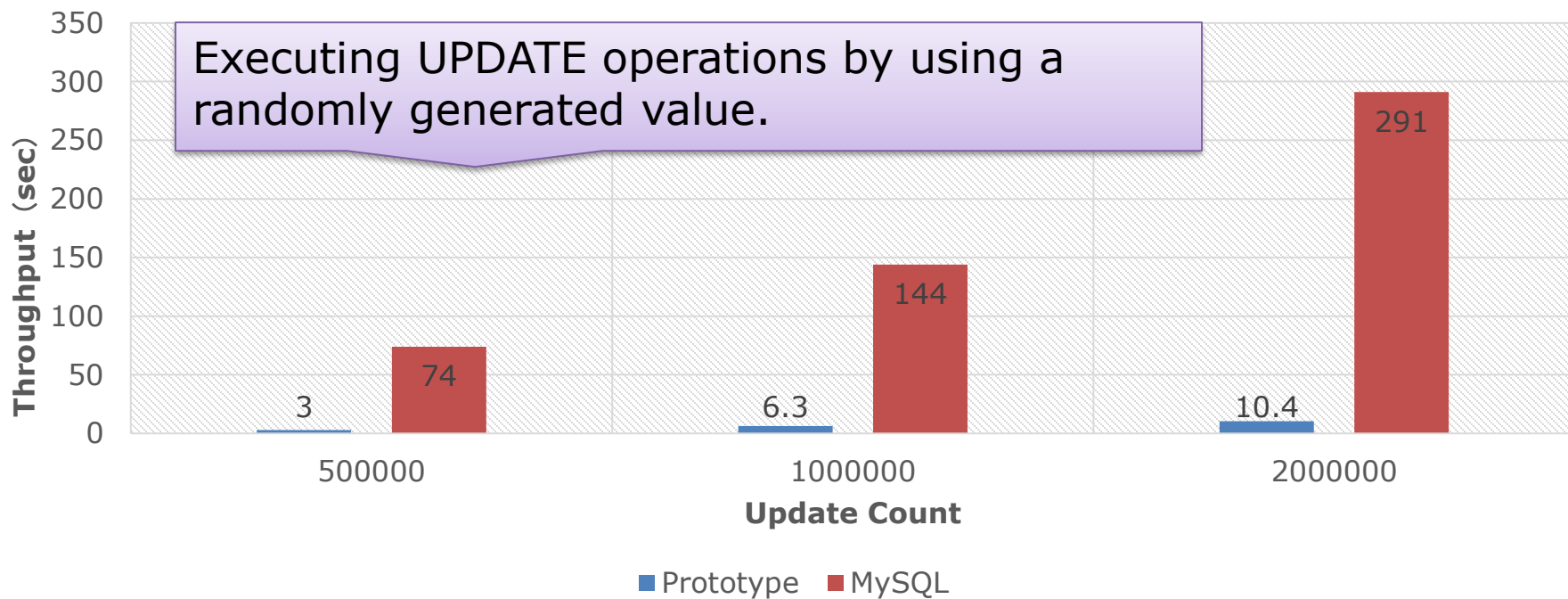
SSD

Every chip can read/write data in parallel

Not so parallel

Highly Parallel

Comparison of the prototype and MySQL: UPDATE command



Theoretical Conclusion

Increase the number
of threads as many
as you can

Implementation!

- Here I introduce a quick implementation method of read-only database engine.

Use Case: Fast Search Engine

SHADOWVERSE
P O R T A L

NEWS CARD LIST DECK BUILDER Sign In

HOME > Card List

Card Name

Class

logical disjunction

ARISA Forestcraft ERIKA Swordcraft ISABELLE Runecraft ROWEN Dragoncraft LUNA Shadowcraft URIAS Bloodcraft ERIS Havencraft Neutral

Cost 1 2 3 4 5 6 7 8 9 10+ Clear

Type ☐ Followers ☐ Spells ☐ Amulets Clear

Rarity ☐ Bronze ☐ Silver ☐ Gold ☐ Legendary Clear

Advanced Search ▲

Search Reset

logical conjunction

Query type detection








SHADOWVERSE
P O R T A L

NEWS CARD LIST DECK BUILDER

HOME > Card List

Card Name

Class

								Clear
Forestcraft	Swordcraft	Runecraft	Dragoncraft	Shadowcraft	Bloodcraft	Havencraft	Neutral	

Cost ☐ 1 ☐ 2 ☐ 3 ☐ 4 ☐ 5 ☐ 6 ☐ 7 ☐ 8 ☐ 9 ☐ 10+

Type ☐ Followers ☐ Spells ☐ Amulets

Rarity ☐ Bronze ☐ Silver ☐ Gold ☐ Legendary

Those are sub-queries corresponds to the immutable queries

Query type detection







SHADOWVERSE
P O R T A L

NEWS CARD LIST DECK BUILDER

HOME > Card List

Card Name

Class

								Clear
Forestcraft	Swordcraft	Runecraft	Dragoncraft	Shadowcraft	Bloodcraft	Havencraft	Neutral	

Cost ☐ 1 ☐ 2 ☐ 3 ☐ 4 ☐ 5 ☐ 6 ☐ 7 ☐ 8 ☐ 9 ☐ 10+

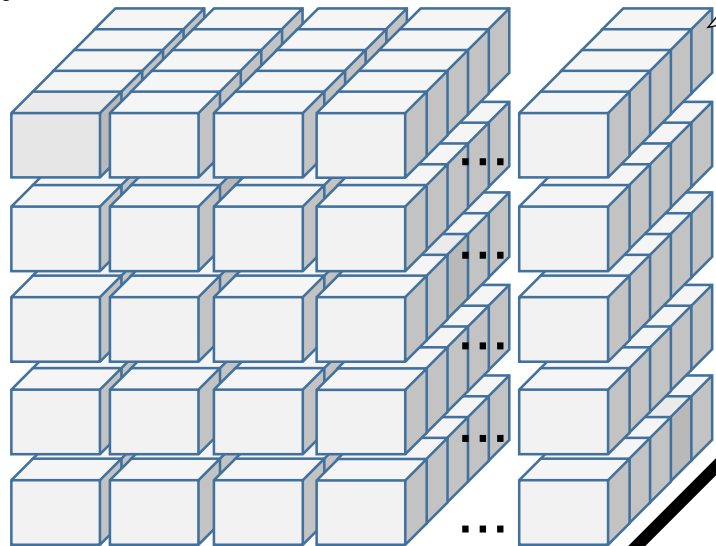
Type ☐ Followers ☐ Spells ☐ Amulets

Rarity ☐ Bronze ☐ Silver ☐ Gold ☐ Legendary

When the target data is fixed, we can execute those sub queries before the runtime.

Data Structure

Attributes



Each block represents an array that stores results of the specific queries applied to the specific attributes.

Query
Results

Queries

Conclusion

- You can implement tiny RDB within a hour.
- Increase the number of threads as many as you can

Recommended Books

- Database Systems: The Complete Book (2nd Edition) 2nd Edition by Hector Garcia-Molina, Jeffrey D. Ullman, Jennifer Widom
- Introduction to Algorithms, 3rd Edition (MIT Press) 3rd Edition by Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, Clifford Stein